## Docker setup

You can start docker pull registry.gitlab.inria.fr/soliman/inf555/td3 now

## Local Search & Constraint Satisfaction Problems

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Thanks to P. Flener, L. Michel and P. Van Hentenryck for inspiration

#### What is "Local Search"

iterative optimization method

 looking for an assignment of variables to values of their domains that minimizes some cost

- local move from solution to neighboring solution
- try to (always or not) decrease the cost of the selected assignment

## How does this relate to Constraint Solving?

Compared to what we have seen up to now:

- only optimization
- requires defining:
  - neighborhood and
  - selection criterion (single state)
  - stopping criterion
- ⇒ **incomplete** (i.e., not optimal) but **low cost** (time and memory)

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Use constraint violations as cost function

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Other are soft constraints that will guide the search Allows us to solve *over-constrained* problems

Neighborhood is defined as changing a single variable assignment

In that framework redundant constraints play two roles:

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#### Violations

#### Violations defined on basic constraints

#### Can be composed:

- $V(c_1 \wedge c_2) = V(c_1) + V(c_2)$
- $V(c_1 \vee c_2) = \min(V(c_1), V(c_2))$
- $V(\neg c) = 1 \min(1, V(c))$

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- $V(c_1 \vee c_2) = \min(V(c_1), V(c_2))$
- $V(\neg c) = 1 \min(1, V(c))$   $\Rightarrow$  not compatible with the above!

For global constraints, one can decompose or not

## N-Queens as a Local Search problem

```
constraint alldifferent(q);
constraint alldifferent([q[i] + i | i in 1..n]);
constraint alldifferent([q[i] - i | i in 1..n]);
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Count violations as the **total** number of identical pairs in an alldifferent constraint

Very dependent on the model! (dual, symmetries, etc.) But the state space does depend too!

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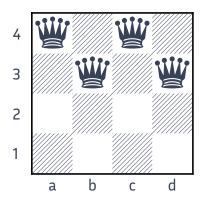
Very dependent on the model! (dual, symmetries, etc.) But the state space does depend too!

Could use the max number of equalities instead of their sum try to guide the search as much as possible

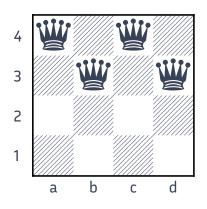
# 

#### alldifferent lists:

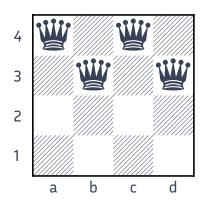
 $\rightarrow$ 



$$\rightarrow$$
[4,3,4,3]  $\searrow$ 

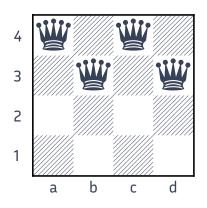


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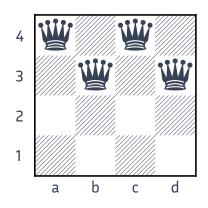
Violations:



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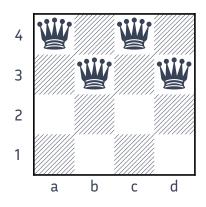
Violations:

2+



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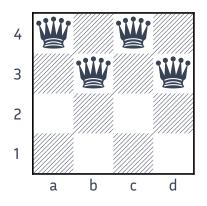


$$\rightarrow$$
[4,3,4,3]  $\searrow$ [4,4,6,6]  $\nearrow$ [4,2,2,0]

Violations:

$$2+2+1=5$$

Neighborhood:



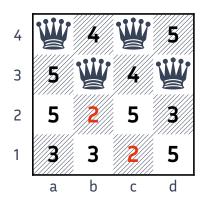
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Neighborhood:

Move **one queen** in its column i.e., change the valuation of a single variable



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#### Neighborhood:

Move **one queen** in its column i.e., change the valuation of a single variable

## Example: Greedy Local Search (aka. Hill-climbing)

Pure exploitation (intensification)

Analog for the discrete case of gradient descent

Select the most improving neighbor (random if multiple bests)

Stop when no improvement found

## Example: Min. Conflict Search (MCS) / Heuristics (MCH)

Original heuristics for CBLS on SAT problems (and still part of GSAT, WalkSAT, etc.)

Implemented by default in the COMET system

Basis of most other heuristics

Select the neighbor (i.e., variable assignment) that minimizes the **number** of violated constraints

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IOW, Hill-Climbing with violations saturated at 1

Local extrema

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Plateaus

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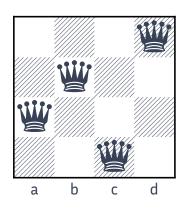
Diagonal ridges (i.e., moves of same cost leading to different extrema)

Local extrema

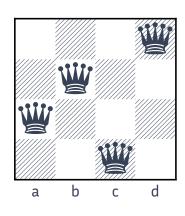
Plateaus

Diagonal ridges (i.e., moves of same cost leading to different extrema)

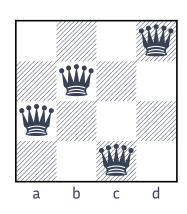
Big neighborhood (might require two steps: **variable** and then **value** selection)



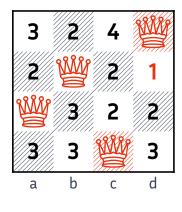
Violation cost =

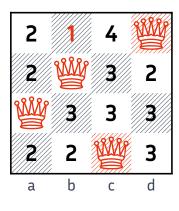


Violation cost = 1 Line-wise costs

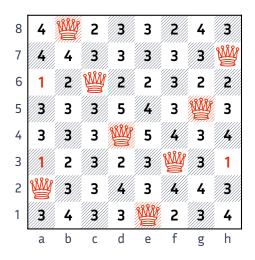


Violation cost = 1 Line-wise costs





Column-wise costs



Line-wise costs

Real local optimum

Necessary to get through a worse solution to get to a global optimum

## Example: Random walk

Pure exploration (diversification)

Select a random neighbor

Remember the best solution found

Stop after a given number of iterations

### Being smarter

⇒ combine a way to escape local extrema with some Hill-climbing

#### Some examples:

diagonal moves (see Practical work session)

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- random restarts (underrated!)

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## Some examples:

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- Tabu search
- random restarts (underrated!)

What if restarts were actually done simultaneously? population-based approaches



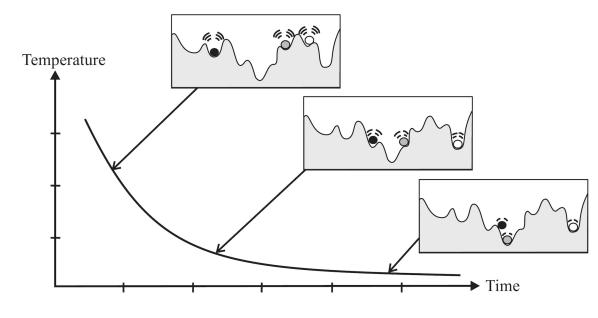
## Simulated annealing

Inspiration from the physics' world (Metropolis-Hastings algorithm for a sampling states of a thermodynamic system, 1953)

Allow some exploration while the temperature of the system is high

Decrease temperature with time (i.e., iterations)

Focus on *exploitation* when the system cools down



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# Practically

### At each step:

- select a random neighbor (no guidance at all...)
- compare its cost with the current cost
- accept it or not depending on the temperature but always accept improving moves
- stop if the move was rejected and the temperature too low (return the best solution found)

Parameters: initial/stopping temperature, cooling regime, acceptance condition

Initial temperature:

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Cooling regime:

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Cooling regime: observed to have almost *no impact* usually  $T_{t+1} = (1-c)T_t$  with a small cooling-rate, e.g., c = 0.01

Acceptance condition:

Initial temperature: allow any move (random-walk)

Ending temperature: would reject most non-improving moves (hill-climbing)

Cooling regime: observed to have almost *no impact* usually  $T_{t+1} = (1-c)T_t$  with a small cooling-rate, e.g., c = 0.01

Acceptance condition: mostly coming from physics consensus:  $\exp(\Delta/T) > r$ ,  $\Delta$  is current cost minus new cost, r is a random variable in [0,1)

## Results

Very often used as a basic LS algorithm

Decent results on the Traveling Salesman Problem not for finding an optimal solution but a *good* solution

On the N-Queens problem...

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On the N-Queens problem... we'll see during the *practical work session* 

Not always easy to fine-tune temperature

## Tabu search

Created by F. Glover in 1986

At its core: Hill-Climbing with some kind of memory forbidding moves

At each step select the **best neighbor** that is **not tabu** 

Stop after a given number of iterations

Tabu moves force diversification

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Tabu moves force diversification in a more *guided* way than random-walks

## Types of memory

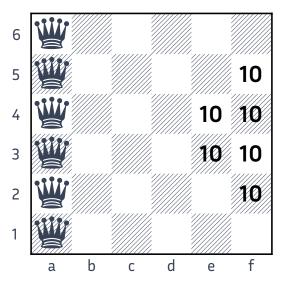
The *Tabu-list* is usually a list of recently visited states, to avoid cycles (classical issue with random diagonal moves)

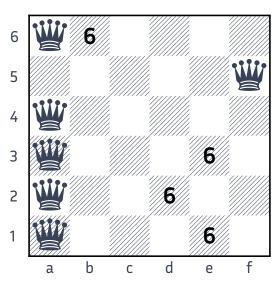
Its length is a sensitive parameter

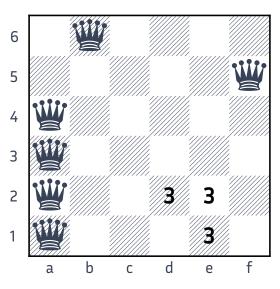
Sometimes not a full state but a feature is stored

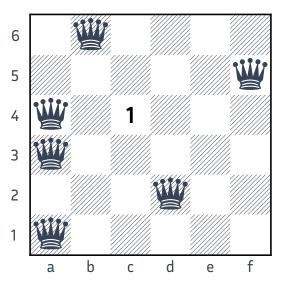
in that case it might be necessary to overcome *tabu* when a better candidate is found

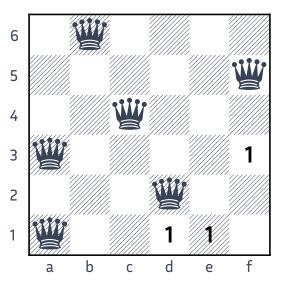
Intermediate-term memory (intensification) and Long-term memory (diversification) rules can be added

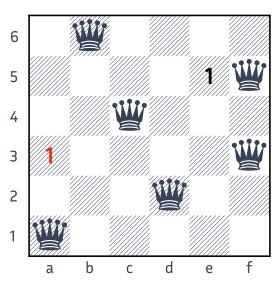


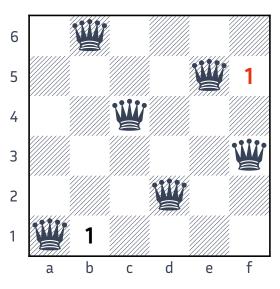


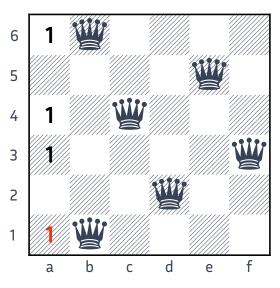


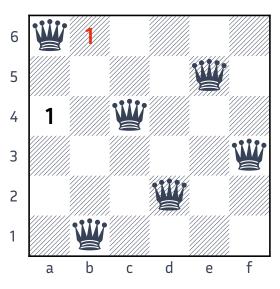


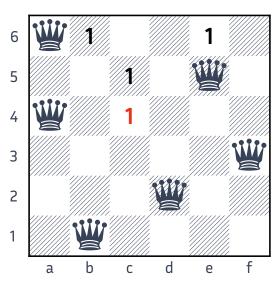


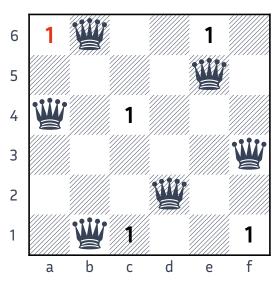


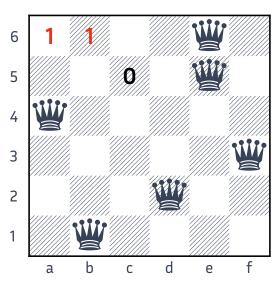


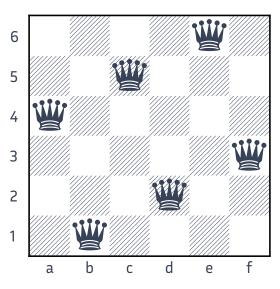












## Results

Same as SA but better/worse

More guided, but more parameters

Good results on TSP

Even trickier to fine-tune (especially complex Tabu structures/rules)

# Population-based approaches

Many versions: Genetic algorithms, Particle Swarm Optimization, Ant-Colony Optimization, etc.

Main idea: instead of restarts, use the information of parallel runs while they are running

Balance between exploration/exploitation, diversification/intensification remains hard

## Conclusion

Incomplete but **efficient** method, even with *simple algorithms* (e.g., Hill-climbing with restarts and diagonal moves)

Used for **hard** problems for which a complete search is not tractable (e.g., Ant Colony on graph problems by C. Solnon)

Or for problems that are **over-constrained** (and can be expressed as optimization) (e.g., MaxSAT)

Can be made generic for CSPs once violations are defined

Remains often tricky to parametrize Does **never prove optimality** 

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