

**MPRI C2-19 Examination, Part II, F. Fages**  
**Model calibration with respect to Temporal Logic**  
**constraints: patterns and solvers**

We consider the fundamental problem of calibrating an ODE model with respect to a temporal logic specification of the behavior of the system, observed in experiments. More precisely, we consider the following grammar of Linear Time Logic formulae with linear constraints over the Reals,  $\text{LTL}(\mathbb{R}_{\text{lin}})$ :

$$\phi ::= c \mid \phi \Rightarrow \psi \mid \phi \wedge \psi \mid \phi \vee \psi \mid \mathbf{X}\phi \mid \mathbf{F}\phi \mid \mathbf{G}\phi \mid \phi \mathbf{U}\psi \mid \phi \mathbf{R}\psi \mid \exists x \phi$$

where  $c$  denotes a linear constraint over molecular concentrations, their first derivative, and a set of variables  $\mathbf{x}$ , including the Time variable.

Timing constraints can be expressed with the time variable and free variables to relate the time of different events. For instance, the formula

$$\mathbf{G}(\text{Time} \leq t_1 \Rightarrow [A] < 1 \wedge \text{Time} \geq t_2 \Rightarrow [A] > 10) \wedge (t_2 - t_1 < 60)$$

expresses that the concentration of molecule  $A$  is always less than 1 up to some time  $t_1$ , always greater than 10 after time  $t_2$ , and the switching time between  $t_1$  and  $t_2$  is less than 60 units of time.

We have seen in the course that the validity domains of the free variables of an  $\text{LTL}(\mathbb{R}_{\text{lin}})$  formula on a finite trace  $(s_0, \dots, s_n)$ , can be computed by finite unions and intersections of the validity domains of the linear constraints, as follows:

$$\begin{aligned} \mathcal{D}_{(s_0, \dots, s_n), \phi} &= \mathcal{D}_{s_0, \phi}, \\ \mathcal{D}_{s_i, c(\mathbf{x})} &= \{\mathbf{v} \in \mathbb{R}^k \mid s_i \models c[\mathbf{v}/\mathbf{x}]\} \text{ for a constraint } c(\mathbf{x}), \\ \mathcal{D}_{s_i, \phi \wedge \psi} &= \mathcal{D}_{s_i, \phi} \cap \mathcal{D}_{s_i, \psi}, \\ \mathcal{D}_{s_i, \phi \vee \psi} &= \mathcal{D}_{s_i, \phi} \cup \mathcal{D}_{s_i, \psi}, \\ \mathcal{D}_{s_i, \mathbf{X}\phi} &= \mathcal{D}_{s_{i+1}, \phi}, \\ \mathcal{D}_{s_i, \mathbf{F}\phi} &= \bigcup_{j=i}^n \mathcal{D}_{s_j, \phi}, \\ \mathcal{D}_{s_i, \mathbf{G}\phi} &= \bigcap_{j=i}^n \mathcal{D}_{s_j, \phi}, \\ \mathcal{D}_{s_i, \phi \mathbf{U}\psi} &= \bigcup_{j=i}^n (\mathcal{D}_{s_j, \psi} \cap \bigcap_{k=i}^{j-1} \mathcal{D}_{s_k, \phi}), \\ \mathcal{D}_{s_i, \exists x \phi} &= \Pi^x \mathcal{D}_{s_i, \phi}, \end{aligned}$$

where the projection operator  $\Pi^x$  returns the domain with all possible values for  $x$ .

This serves to compute a continuous violation degree for a  $\text{LTL}(\mathbb{R}_{\text{lin}})$  formula with objective values  $v^*$  for the variables as the euclidian distance between the validity domain and the objective values:

$$vd_{s_i, \phi}^{v^*} = \text{dist}(\mathcal{D}_{s_i, \phi}, v^*)$$

Furthermore, a “robustness” degree can be estimated as the interior distance to the borders of the validity domain, i.e. as the distance to the complement of the validity domain:

$$ro_{s_i, \phi}^{v^*} = \text{dist}(\overline{\mathcal{D}_{s_i, \phi}}, v^*).$$

We are interested in finding more efficient algorithms for computing these values for some common patterns of  $\text{LTL}(\mathbb{R}_{\text{lin}})$  formulae. For instance, for the constraint

$$A \geq v$$

where  $v$  is a free variable, we have

$$\begin{aligned}\mathcal{D}_{s_i, A \geq v} &= ] -\infty, \text{val}_{s_i}(A)] \\ \text{vd}_{s_i, A \geq v}^{v^*} &= \max(0, v^* - \text{val}_{s_i}(A)) \\ \text{ro}_{s_i, A \geq v}^{v^*} &= \max(0, \text{val}_{s_i}(A) - v^*)\end{aligned}$$

where  $\text{val}_{s_i}(A)$  is the value of  $A$  in state  $s_i$ ,  
For the reachability formula

$$\mathbf{F}(A \geq v)$$

we have

$$\begin{aligned}\mathcal{D}_{s_i, \mathbf{F}(A \geq v)} &= ] -\infty, \text{max}_{s_i}(A)] \\ \text{vd}_{s_i, \mathbf{F}(A \geq v)}^{v^*} &= \max(0, v^* - \text{max}_{s_i}(A)) \\ \text{ro}_{s_i, \mathbf{F}(A \geq v)}^{v^*} &= \max(0, \text{max}_{s_i}(A) - v^*)\end{aligned}$$

where  $\text{max}_{s_i}(A)$  is the maximum value of  $A$  over the trace  $(s_i, \dots, s_n)$ , which can be computed in one pass over the trace, in time  $O(n - i)$ .

### Give similar expressions for the constraint $A \leq v$

$$\begin{aligned}\mathcal{D}_{s_i, A \leq v} &= [\text{val}_{s_i}(A), +\infty[ \\ \text{vd}_{s_i, A \leq v}^{v^*} &= \max(0, \text{val}_{s_i}(A) - v^*) \\ \text{ro}_{s_i, A \leq v}^{v^*} &= \max(0, v^* - \text{val}_{s_i}(A))\end{aligned}$$

### Give similar expressions and justify them for the formula $\mathbf{F}(A \leq v)$

$$\mathcal{D}_{s_i, \mathbf{F}(A \leq v)} = \bigcup_{j=i}^n \mathcal{D}_{s_j, A \leq v} = \bigcup_{j=i}^n [\text{val}_{s_j}(A), +\infty[ = [\text{min}_{s_i}(A), +\infty[$$

$$\text{vd}_{s_i, \mathbf{F}(A \leq v)}^{v^*} = \text{dist}([\text{max}_{s_i}(A), +\infty[, v^*) = \max(0, \text{min}_{s_i}(A) - v^*)$$

$$\text{ro}_{s_i, \mathbf{F}(A \leq v)}^{v^*} = \text{dist}(] -\infty, \text{max}_{s_i}(A)[, v^*) = \max(0, v^* - \text{min}_{s_i}(A))$$

**Give similar expressions for the formula  $\mathbf{G}(A \leq v)$**

$$\mathcal{D}_{s_i, \mathbf{G}(A \leq v)} = \bigcap_{j=i}^n \mathcal{D}_{s_j, A \leq v} = \bigcap_{j=i}^n [val_{s_j}(A), +\infty[ = [max_{s_i}(A), +\infty[$$

$$vd_{s_i, \mathbf{G}(A \leq v)}^{v^*} = dist([max_{s_i}(A), +\infty[, v^*) = max(0, max_{s_i}(A) - v^*)$$

$$ro_{s_i, \mathbf{G}(A \leq v)}^{v^*} = dist(] - \infty, max_{s_i}(A)[, v^*) = max(0, v^* - max_{s_i}(A))$$

**Give similar expressions for the amplitude constraint  $\exists v (\mathbf{F}(A \leq v) \wedge \mathbf{F}(A \geq v+a))$**

The vector of variables by the tuple  $(v, a)$  in this order.

$$\mathcal{D}_{s_i, \mathbf{F}(A \leq v) \wedge \mathbf{F}(A \geq v+a)} = \{(v, a) \in \mathbf{R}^2 \mid min_{s_i} \leq v \wedge max_{s_i} \geq v+a\}$$

$$\begin{aligned} \mathcal{D}_{s_i, \exists v (\mathbf{F}(A \leq v) \wedge \mathbf{F}(A \geq v+a))} &= \Pi^v(\{(v, a) \in \mathbf{R}^2 \mid min_{s_i} \leq v \wedge max_{s_i} \geq v+a\}) \\ &= (\mathbf{R}, ] - \infty, max_{s_i}(A) - min_{s_i}(A)]) \end{aligned}$$

$$vd_{s_i, \exists v (\mathbf{F}(A \leq v) \wedge \mathbf{F}(A \geq v+a))}^{a^*} = max(0, a^* - max_{s_i}(A) + min_{s_i}(A))$$

$$ro_{s_i, \exists v (\mathbf{F}(A \leq v) \wedge \mathbf{F}(A \geq v+a))}^{a^*} = max(0, max_{s_i}(A) - min_{s_i}(A) - a^*)$$