

From Reaction Models to Influence Models and Back

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1. Theory of abstract interpretation
2. Syntactical domain of SBML reaction rules
3. Hierarchy of Stochastic to Boolean semantics
4. Activation/inhibition influences abstracted from the syntax of reaction rules
5. Activation/inhibition influences abstracted from the differential semantics of reaction rules
6. Conclusion

1. Abstraction in Systems Biology

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These perspectives can be reconciled by organizing models into **hierarchies of abstractions.**

“To understand a system is not to know everything about it but to know abstraction levels that are sufficient for answering questions about it”

The Theory of Abstract Interpretation

In this setting [Cousot Cousot 77], a **domain** is a lattice $\mathcal{D}(\sqsubseteq, \perp, \top, \sqcup, \sqcap)$ where \sqsubseteq is the “information loss” ordering.

Often just a power-set $\mathcal{P}(\mathcal{S})(\subseteq, \emptyset, \mathcal{S}, \cup, \cap)$ ordered by set inclusion.

The Theory of Abstract Interpretation

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A **Galois connection** $\mathcal{C} \rightarrow_{\alpha} \mathcal{A}$ between two lattices \mathcal{C} and \mathcal{A} is defined by two abstraction and concretization functions $\alpha : \mathcal{C} \rightarrow \mathcal{A}$ and $\gamma : \mathcal{A} \rightarrow \mathcal{C}$ that are **monotonic**:

- $\forall x, y \in \mathcal{C} \ x \sqsubseteq_{\mathcal{C}} y \Rightarrow \alpha(x) \sqsubseteq_{\mathcal{A}} \alpha(y),$
- $\forall x, y \in \mathcal{A} \ x \sqsubseteq_{\mathcal{A}} y \Rightarrow \gamma(x) \sqsubseteq_{\mathcal{C}} \gamma(y),$

and are **adjoint**:

- $\forall c \in \mathcal{C}, \forall y \in \mathcal{A} : c \sqsubseteq_{\mathcal{C}} \gamma(y) \Leftrightarrow \alpha(c) \sqsubseteq_{\mathcal{A}} y.$

If $\gamma \circ \alpha$ is the identity, the abstraction α loses no information, and \mathcal{C} and \mathcal{A} are isomorphic from the information standpoint (although α may be not onto and γ not one-to-one).

Properties of Galois Connections

1. $\gamma \circ \alpha$ is extensive (i.e. $x \sqsubseteq_{\mathcal{C}} \gamma \circ \alpha(x)$) and represents the information lost by the abstraction;
2. $\alpha \circ \gamma$ is contracting (i.e. $\alpha \circ \gamma(y) \sqsubseteq_{\mathcal{A}} y$);
3. $\gamma \circ \alpha$ is the identity *iff* γ is onto *iff* α is one-to-one.
4. α preserves \sqcup , and γ preserves \sqcap ;
5. $\gamma(a) = \max \alpha^{-1}(\downarrow a) = \sqcup \alpha^{-1}(\downarrow a)$
6. $\alpha(c) = \min \gamma^{-1}(\uparrow c) = \sqcap \gamma^{-1}(\uparrow c)$

where $\downarrow a = \{b \mid b \sqsubseteq a\}$ and $\uparrow a = \{b \mid a \sqsubseteq b\}$.

It is equivalent in the definition of Galois connections to replace the condition of adjointness by conditions 1 and 2,

or by condition 5 which also entails the monotonicity of γ .

Systems Biology Markup Language SBML Models

Formally, the concrete domain of reaction models is the powerset of all possible reaction rules ordered by set inclusion :

Def. 1 *Given a finite set \mathcal{M} of molecule names, the universe of reactions is the set of rules*

$$\mathcal{R} = \{e \text{ for } S \Rightarrow S' \mid \begin{array}{l} e \text{ is a kinetic expression,} \\ \text{and } S \text{ and } S' \text{ are solutions of molecules in } \mathcal{M} \end{array}\}.$$

The domain of SBML reaction models is $\mathcal{C}_{\mathcal{R}} = (\mathcal{P}(\mathcal{R}), \subseteq)$.

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In the SBML exchange format, no semantics are defined.

In BIOCHAM, three semantics are considered:

1. boolean : non-deterministic asynchronous transition system
2. differential : ODE (or hybrid system)
3. stochastic : continuous time Markov chain.

Stochastic Semantics Domain

Def. 3 Let a *discrete state* be a vector of integers of dimension $|\mathcal{M}|$. The universe \mathcal{S} of *stochastic transitions* is the set of triplets (S_i, S_j, τ) where S_i and S_j are discrete states and $\tau \in \mathbb{R}^+$ is a weight for the transition.

The domain of stochastic transitions is $\mathcal{D}_{\mathcal{S}} = (\mathcal{P}(\mathcal{S}), \subseteq)$.

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An element s of the domain precisely defines a Markov chain, where the probability p_{ij} of transition from state S_i to S_j is obtained by normalizing

the reaction rate $\tau_{i,j} = \sum_{(S_i, S_j, \tau) \in s} \tau$ in $p_{ij} = \frac{\tau_{ij}}{\sum_{(S_i, S_k, \tau) \in s} \tau}$

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Proposition 6 Let $\alpha_{\mathcal{R}\mathcal{S}} : \mathcal{C}_{\mathcal{R}} \rightarrow \mathcal{D}_{\mathcal{S}}$ be the function associating to a reaction model the state transition graph labelled with the $\tau_{i,j}$'s. Let

$\gamma_{\mathcal{R}\mathcal{S}}(s) = \cup \alpha_{\mathcal{R}\mathcal{S}}^{-1}(\downarrow s)$. $\mathcal{C}_{\mathcal{R}} \xrightarrow{\alpha_{\mathcal{R}\mathcal{S}}} \mathcal{D}_{\mathcal{S}} \xleftarrow{\gamma_{\mathcal{R}\mathcal{S}}} \mathcal{C}_{\mathcal{R}}$ is a Galois connection.

Remark that $\alpha_{\mathcal{R}\mathcal{S}}$ is not one-to-one nor onto

Discrete Semantics

Def. 7 *The universe \mathcal{D} of discrete transitions is the set of pairs of discrete states. The domain of discrete transitions is $\mathcal{D}_{\mathcal{D}} = (\mathcal{P}(\mathcal{D}), \subseteq)$.*

The discrete semantics is the classical Petri net semantics of reaction models [RML93ismb,SHK06bmcbi,Chaouiya07bioinfo,GHL07cmsb].

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Classical Petri net analysis tools can be used for the analysis of reaction models at this abstraction level. For instance, the elementary mode analysis of metabolic networks [SPM02bioinfo] has been shown in [ZS03insilicobio] to be equivalent to the classical analysis of Petri nets by T-invariants.

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Proposition 10 *Let $\alpha_{\mathcal{S}\mathcal{D}} : \mathcal{D}_{\mathcal{S}} \rightarrow \mathcal{D}_{\mathcal{D}}$ be the function associating to a set of stochastic transitions the discrete transitions obtained by projection on the two first components, and $\gamma_{\mathcal{S}\mathcal{D}}(d) = \cup \alpha_{\mathcal{S}\mathcal{D}}^{-1}(\downarrow d)$. $\mathcal{D}_{\mathcal{S}} \xrightarrow{\alpha_{\mathcal{S}\mathcal{D}}} \mathcal{D}_{\mathcal{D}} \xleftarrow{\gamma_{\mathcal{S}\mathcal{D}}}$ is a Galois connection.*

Remark that $\alpha_{\mathcal{S}\mathcal{D}}$ is onto, but not one-to-one

Boolean Semantics

Def. 11 Let a *boolean state* be a vector of booleans of dimension $|\mathcal{M}|$ indicating the presence of each molecule in the state. The universe \mathcal{B} of *boolean transitions* is the set of pairs of boolean states.

The domain of boolean transitions is $\mathcal{D}_{\mathcal{B}} = (\mathcal{P}(\mathcal{B}), \subseteq)$.

Boolean Semantics

Def. 12 Let a *boolean state* be a vector of booleans of dimension $|\mathcal{M}|$ indicating the presence of each molecule in the state. The universe \mathcal{B} of *boolean transitions* is the set of pairs of boolean states.

The domain of boolean transitions is $\mathcal{D}_{\mathcal{B}} = (\mathcal{P}(\mathcal{B}), \subseteq)$.

Let $\alpha_{\mathcal{N}\mathcal{B}} : \mathbb{N}^{|\mathcal{M}|} \rightarrow \mathbb{B}^{|\mathcal{M}|}$ be the **zero/non-zero abstraction** (or threshold abstraction) from the integers to the booleans, and its pointwise extension from discrete states to boolean states.

Proposition 13 Let $\alpha_{\mathcal{D}\mathcal{B}} : \mathcal{D}_{\mathcal{D}} \rightarrow \mathcal{D}_{\mathcal{B}}$ be the set extension of $\alpha_{\mathcal{N}\mathcal{B}}$. Let $\gamma_{\mathcal{D}\mathcal{B}}(b) = \cup \alpha_{\mathcal{D}\mathcal{B}}^{-1}(\downarrow b)$. $\mathcal{D}_{\mathcal{D}} \xrightarrow{\alpha_{\mathcal{D}\mathcal{B}}} \mathcal{D}_{\mathcal{B}} \xleftarrow{\gamma_{\mathcal{D}\mathcal{B}}} \mathcal{D}_{\mathcal{D}}$ is a Galois connection.

BIOCHAM boolean semantics S_{BB} with the combinatorics of reactant consumption

Theorem 14 For any reaction model R , $\alpha_{\mathcal{D}\mathcal{B}}(\alpha_{\mathcal{S}\mathcal{D}}(\alpha_{\mathcal{R}\mathcal{S}}(R))) \subseteq S_{BB}$.

Differential Semantics ?

The **differential semantics** of reaction models interprets a set of reaction rules $\{e_i \text{ for } S_i \Rightarrow S'_i\}_{i=1,\dots,n}$ over molecular concentration variables $\{x_1, \dots, x_m\}$, by the following system of Ordinary Differential Equations (ODE):

$$dx_k/dt = \sum_{i=1}^n r_i(x_k) * e_i - \sum_{j=1}^n l_j(x_k) * e_j$$

where we recall that $r_i(x_k)$ (resp. l_i) is the stoichiometric coefficient of x_k in the right (resp. left) member of rule i .

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- synchronous semantics (evolution of variables in parallel)
- deterministic semantics (average behavior)
- not compatible with the rule set inclusion ordering
- infinite number of molecules
- infinitesimal time steps

Influence Graph Abstraction from the Reaction Rule Syntax

$$\mathcal{A}_I = \mathcal{P}(\{A \text{ activates } B \mid A, B \in \mathcal{M}\} \cup \{A \text{ inhibits } B \mid A, B \in \mathcal{M}\}).$$

The **influence graph of a reaction model** is defined by $\alpha_{\mathcal{R}I} : \mathcal{C}_{\mathcal{R}} \rightarrow \mathcal{A}_I$

$$\begin{aligned} \alpha_{\mathcal{R}I}(x) = & \{A \text{ inhibits } B \mid \exists(e_i \text{ for } S_i \Rightarrow S'_i) \in x, \\ & l_i(A) > 0 \text{ and } r_i(B) - l_i(B) < 0\} \\ & \cup \{A \text{ activates } B \mid \exists(e_i \text{ for } S_i \Rightarrow S'_i) \in x, \\ & l_i(A) > 0 \text{ and } r_i(B) - l_i(B) > 0\} \end{aligned}$$

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$$\alpha_{\mathcal{R}I}(\{A + B \Rightarrow C\}) = \{ \begin{array}{l} A \text{ inhibits } B, A \text{ inhibits } A, B \text{ inhibits } A, \\ B \text{ inhibits } B, A \text{ activates } C, B \text{ activates } C \end{array} \}$$

$$\alpha_{\mathcal{R}I}(\{A = [C] \Rightarrow B\}) = \{ C \text{ inhibits } A, A \text{ inhibits } A, A \text{ activates } B, C \text{ activates } B \}$$

$$\alpha_{\mathcal{R}I}(\{A = [B] \Rightarrow -\}) = \{ B \text{ inhibits } A, A \text{ inhibits } A \}$$

$$\alpha_{\mathcal{R}I}(\{- = [B] \Rightarrow A\}) = \{ B \text{ activates } A \}$$

Influence Graph Abstraction from the Reaction Rule Syntax

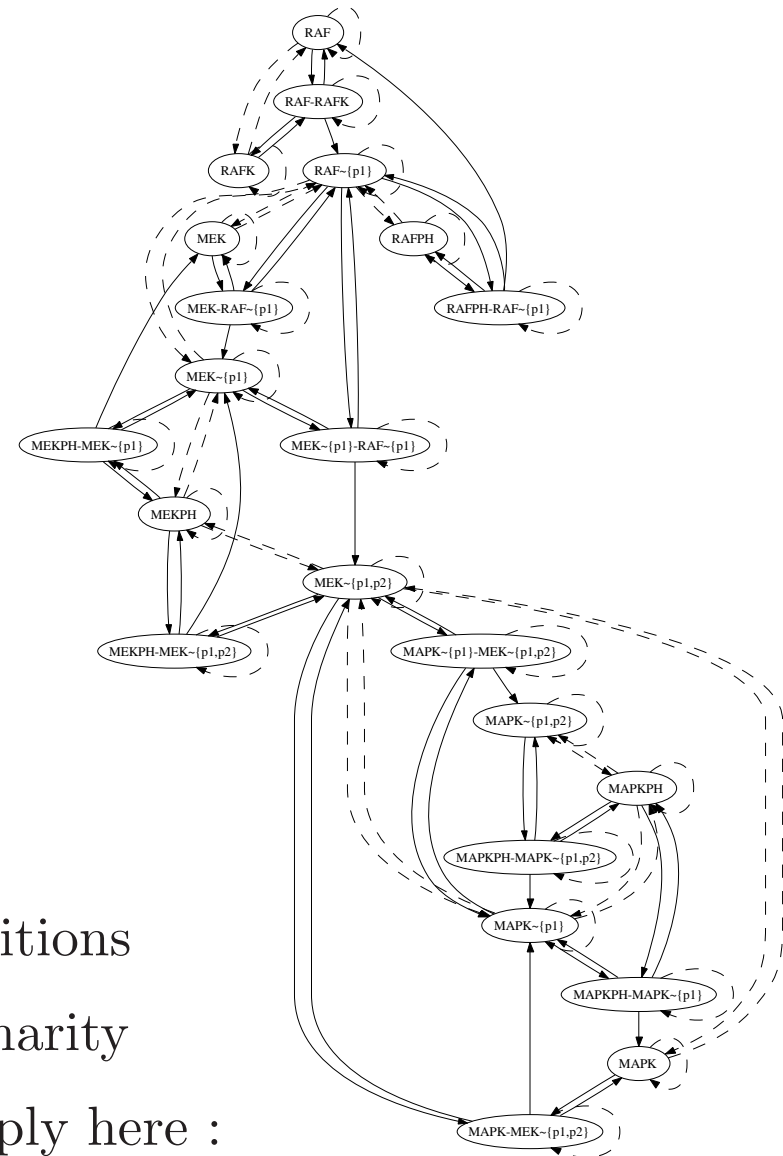
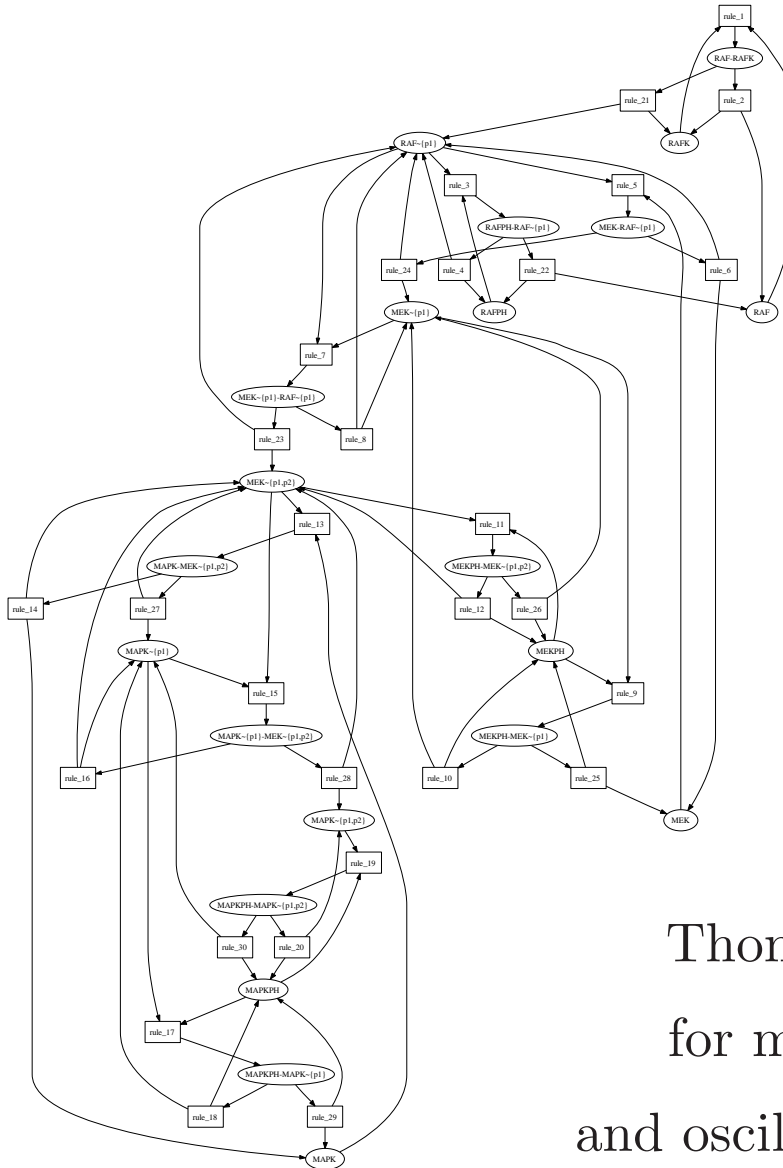
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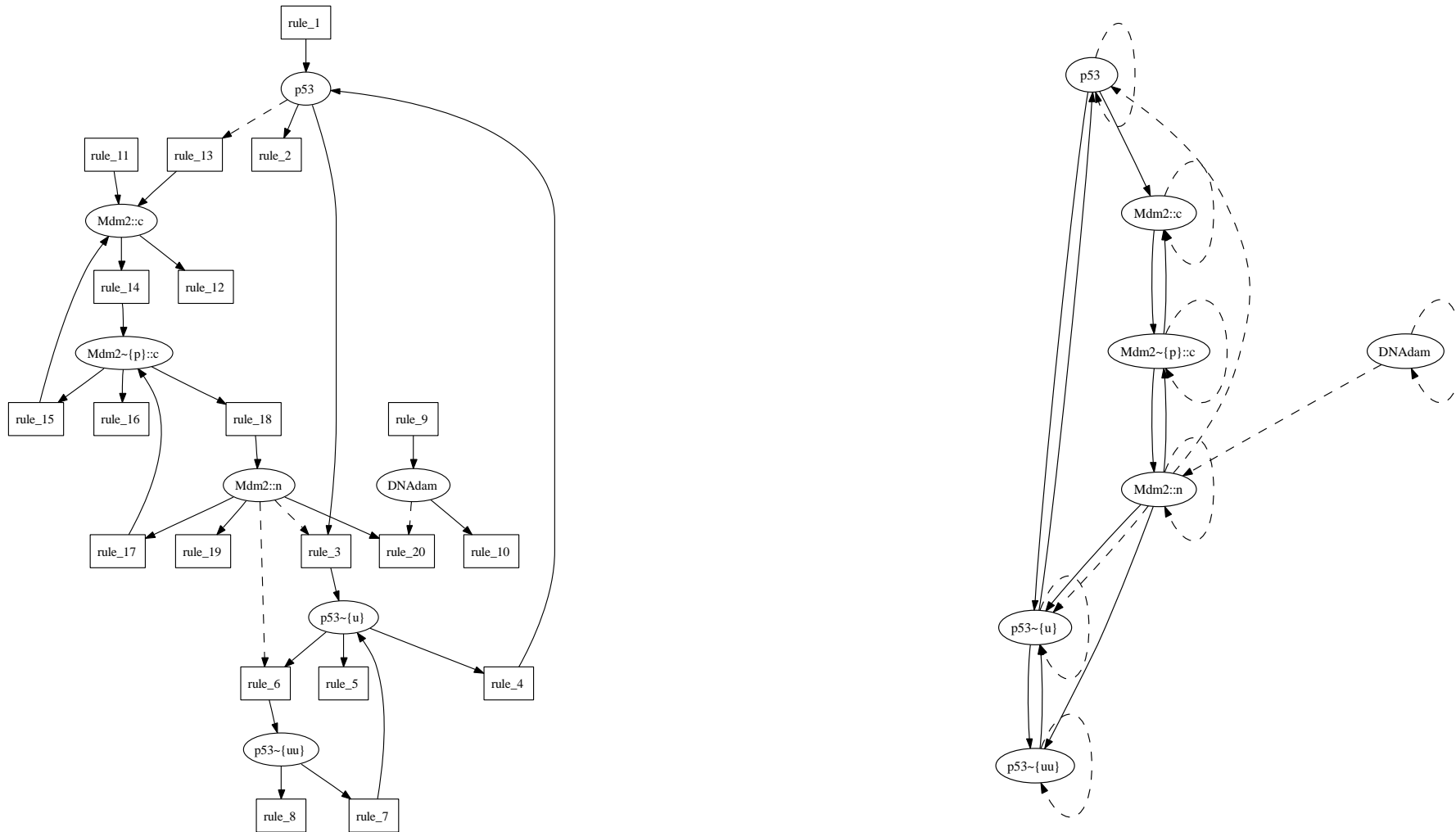
Proposition 15 *Let $\gamma_{\mathcal{R}I}(f) = \cup \alpha_{\mathcal{R}I}^{-1}(\downarrow f)$, $\mathcal{C}_{\mathcal{R}} \xrightarrow{\alpha_{\mathcal{R}I}} \mathcal{A}_I \xleftarrow{\gamma_{\mathcal{R}I}}$ is a Galois connection.*

MAPK model: Reaction Graph \rightarrow_{α} Influence Graph



Thomas's conditions
for multistationarity
and oscillations apply here :

P53-Mdm2: Reaction Graph \rightarrow_{α} Influence Graph



Inhibitions hidden in the kinetic expressions are missed

Influence Graph Abstraction from the Differential Semantics

Let us denote by β the mapping from $\mathcal{C}_{\mathcal{R}}$ to $\mathcal{D}_{\mathcal{J}}$ that extracts \dot{x}_k and hence the Jacobian from the kinetic expressions in the reaction rules.

Def. 16 *The differential influence abstraction $\alpha_{\mathcal{J}\mathcal{I}} : \mathcal{D}_{\mathcal{J}} \rightarrow \mathcal{A}_{\mathcal{I}}$ is the function*

$$\alpha_{\mathcal{J}\mathcal{I}}(x) = \{A \text{ activates } B \mid \partial \dot{x}_B / \partial x_A > 0 \text{ in some point of the space}\} \\ \cup \{A \text{ inhibits } B \mid \partial \dot{x}_B / \partial x_A < 0 \text{ in some point of the space}\}$$

defined purely from the kinetic expressions... compatibility with the rules ?

Monotonic Kinetics

Def. 17 A kinetic expression e_i is *monotonic* w.r.t. a reaction model x iff for all molecules x_k we have

1. for all points of the space $\partial e_i / \partial x_k \geq 0$
2. if there exists a point in the space s.t. $\partial e_i / \partial x_k > 0$ then $l_i(x_k) > 0$

The mass action law kinetics, $e_i = k * \prod x_i^{l_i}$, are monotonic

Hill's kinetics (and Michaelis-Menten kinetics when $n = 1$)

$e_i = V_m * x_s^n / (K_m + x_s^n)$ where $V_m = k * (x_e + x_e * x_s / K_m)$ for an enzymatic reaction $x_s = [x_e] \Rightarrow x_p$, are also monotonic.

Theorem 18 For any reaction model x with monotonic kinetics,

$$\alpha_{JI} \circ \beta(x) \subseteq \alpha_{RI}(x).$$

Precise Kinetics

Def. 19 A kinetic expression e_i is *precise* w.r.t. a reaction model x iff for all molecules x_k we have

1. for all points of the space $\partial e_i / \partial x_k \geq 0$
2. there exists a point in the space s.t. $\partial e_i / \partial x_k > 0$ iff $l_i(x_k) > 0$

Note that *precise implies monotonic*.

Proposition 20 Mass action law, Michaelis Menten, and Hill kinetics are *precise*.

Theorem 21 If x has precise kinetics and no molecule is at the same time an activator and an inhibitor of the same target molecule, then

$$\alpha_{RI}(x) = \alpha_{JI} \circ \beta(x).$$

Conclusion

- The algebraic setting of abstract interpretation applies to
 1. the organization of the syntactical, stochastic, discrete, and boolean semantics of reaction models into a hierarchy of semantics
 2. the mathematical definition of activation/inhibition influence graphs either from the syntax of the rules, or from the differential semantics
- the differential semantics does not belong to that hierarchy of semantics (with the rule set inclusion ordering)
- *thm 1* for monotonic kinetics (like all standard kinetics), the syntactical influences over-approximate the differential influences
- *thm 2* for precise kinetics with no both activation and inhibition influences, the syntactical and differential influences are the same
- the (easily computable) syntactical influences reveal hidden negative feedback in purely forward reaction cascades like MAPK [Sepulchre 07]